

## Greedy Algorithm

### Principle

- At each step, a choice is made, the one that seems the best at that moment
- Builds a solution step by step
  - without revisiting previous decisions
  - by making at each step the choice that seems the best
  - hoping to achieve a global optimal result
- Greedy approach
  - depending on the problem, no guarantee of optimality (greedy heuristic)
  - low cost (compared to exhaustive enumeration)
  - intuitive choice

## Local Search

### Principle

- Start from an initial solution
- At each step, modify the solution
  - trying to improve the value of the objective function
  - hoping to achieve the global optimum
- Local approach
  - depending on the problem, no guarantee of optimality (heuristic)
  - low cost

### Initial solution

- “Empty” solution
- Random solution
- Solution from a greedy algorithm

## Local Search

### Principle

- Start from an initial solution
- At each step, modify the solution
  - trying to improve the value of the objective function
  - hoping to achieve the global optimum
- Local approach
  - depending on the problem, no guarantee of optimality (heuristic)
  - low cost

### Modifications

- Modify the value of a variable
- Swap the values of two variables

## Problem Solving Local Search

Marie Pelleau

marie.pelleau@univ-cotedazur.fr

## Knapsack Problem

### Description

You have:

- A backpack with a weight limit
- A set of objects, each object  $o_i$  has
  - A weight:  $w_i$
  - A value:  $v_i$

Which objects should be taken to maximize the total value carried while respecting the weight constraint?

- The total value of the selected objects is maximized
- The total weight of the selected objects is less than or equal to the backpack's weight limit

## Knapsack Problem

### Variables

- We associate each item with a 0-1 variable (it only takes values 0 or 1)
- It is a membership variable for the backpack
- If the item is taken, the variable is 1, otherwise it is 0

### Model

- The value and weight of an item are given data, so for item  $o_i$ , we have the value  $v_i$  and the weight  $w_i$
- The membership variable for the backpack is  $x_i$
- The maximum weight of the backpack is  $W$

## Knapsack Problem

### Constraints

- $\max \sum_{i=1}^n v_i x_i$  the objective
- $\sum_{i=1}^n w_i x_i \leq W$  sum of weights less than or equal to the maximum weight

## Knapsack

### Initial solution

- "Empty" solution: empty knapsack  $\Rightarrow$  objective function 0
- Random solution : random backpack  $\Rightarrow$  must be verified as a solution
- Solution of a greedy algorithm

### Modifications

- Add an item to the backpack  $\Rightarrow$  if max capacity is not exceeded
- Deletes an item from the backpack

## Hitting-set: Set cover

## Description

- A switch is connected to some bulbs
- When a switch is pressed, all connected bulbs are turned on
- **Question:** What is the minimum number of switches needed to turn on all bulbs?

## Hitting-set: Set cover

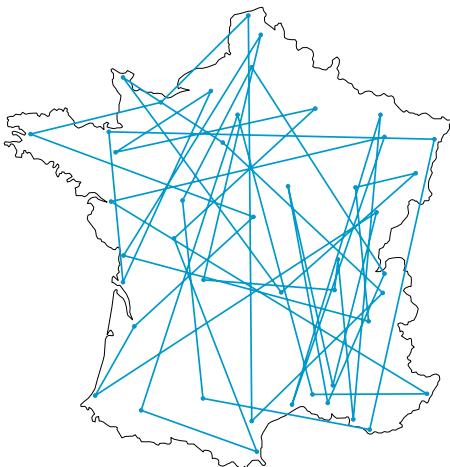
## Initial solution

- "Empty" solution: all switches on  $\Rightarrow$  objective function number of switches
- Random solution: random switch positions  $\Rightarrow$  must be verified as a solution
- Solution of a greedy algorithm

## Modifications

- Turn on a switch
- Turns off a switch  $\Rightarrow$  if all bulbs remain on

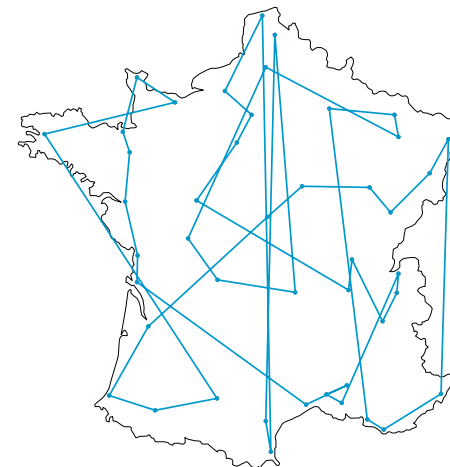
## TSP



## Solution initiale

- Cities in alphabetical order

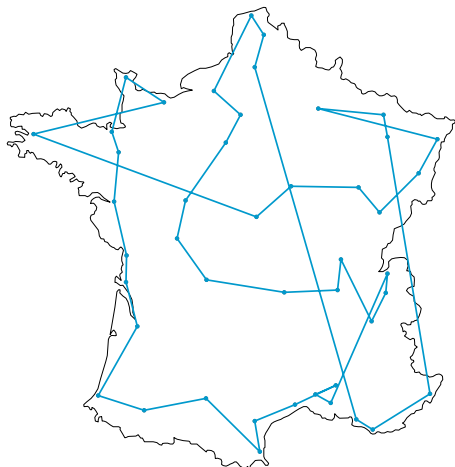
## TSP



## Solution initiale

- Cities in alphabetical order
- Cities in random order

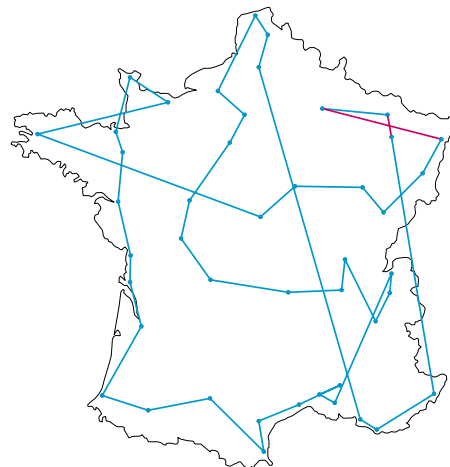
## TSP



## Solution initiale

- Cities in alphabetical order
- Cities in random order
- Solution of a greedy algorithm

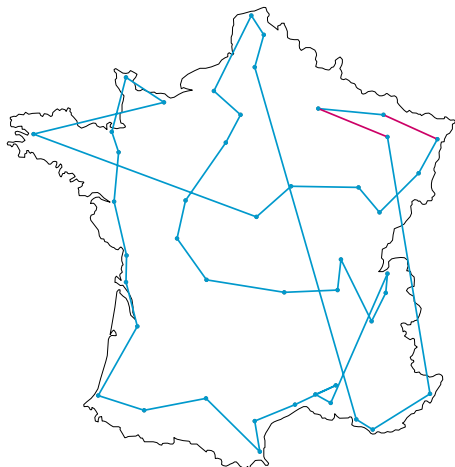
## TSP



## Modifications

- $k$ -opt
  - $k = 2$
  - $k = 3$

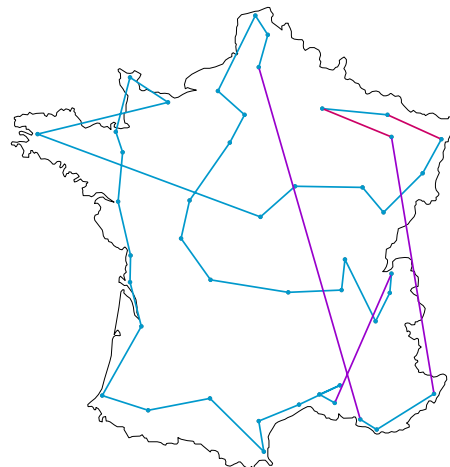
## TSP



## Modifications

- $k$ -opt
  - $k = 2$
  - $k = 3$

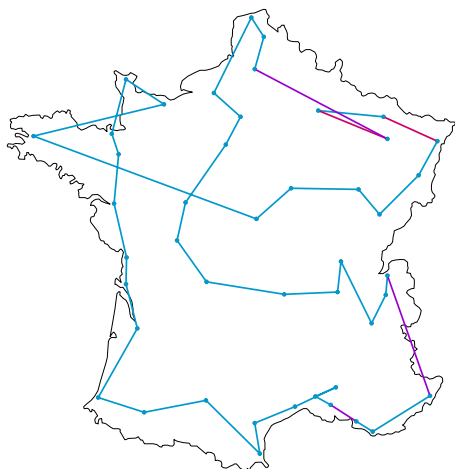
## TSP



## Modifications

- $k$ -opt
  - $k = 2$
  - $k = 3$

## TSP



## Modifications

- $k$ -opt
  - $k = 2$
  - $k = 3$

## Local search

## Principle

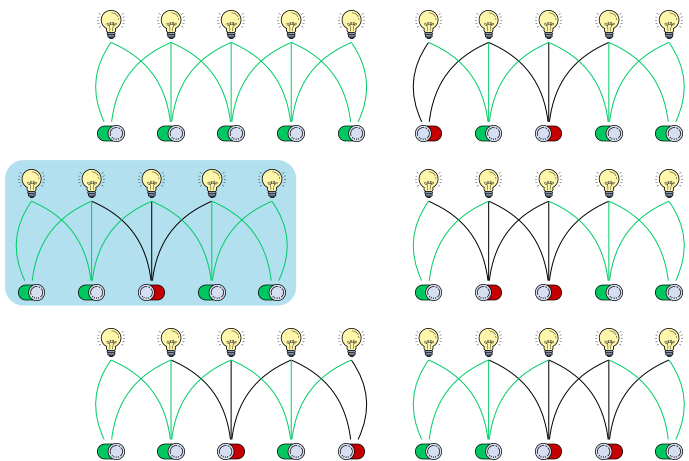
- We start with an initial solution
- At each step, we modify the solution  $\Rightarrow$  notion of neighborhood

## Neighborhood

For a solution, the set of solutions with one modification

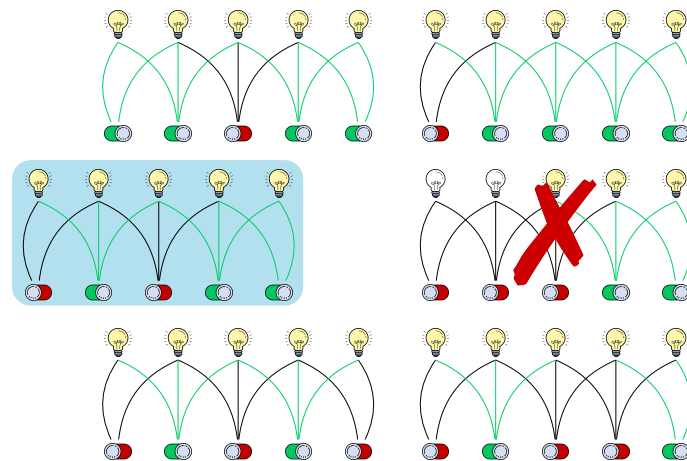
## Hitting-set: Set cover

## Neighborhood



## Hitting-set: Set cover

## Neighborhood



## Local search

### Principle

- We start with an initial solution
- At each step, we modify the solution  $\Rightarrow$  notion of neighborhood

### Which neighbor to choose?

- Randomly
- The best
- One of the best

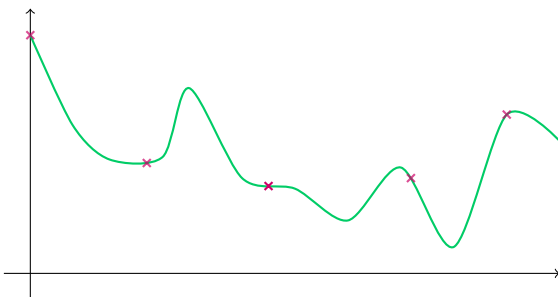
## Contents

- 1 Random walk
- 2 Gradient descent
- 3 Restarts
- 4 Tabu Search
- 5 Constraint Based Local Search

## Random walk

### Principle

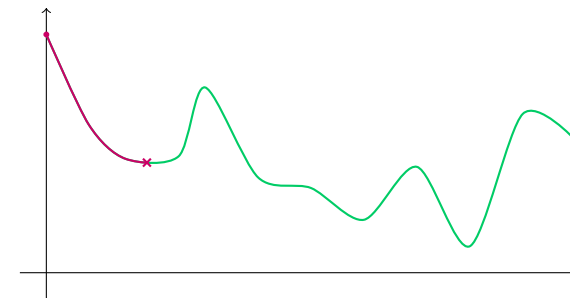
- We start with an initial solution
- At each step, the solution is **randomly** modified



## Gradient descent

### Principle

- We start with an initial solution
- At each step, we move towards a solution in the neighborhood **strictly improving** the objective



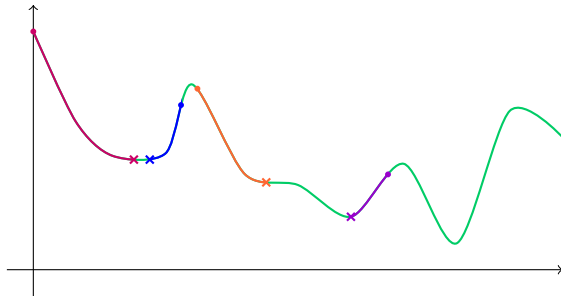
### Drawbacks

You can get stuck in local minima

## Gradient descent

### Principle

- We start with an initial solution
- At each step, we move towards a solution in the neighborhood **strictly improving** the objective



### Restarts

Start again from another solution

## Local search

### Restarts

- Random solution
- “Empty” solution, in which a certain percentage of variables is fixed as in the best solution found so far
  - 5%, 10%, 20%

**Large Neighborhood Search** (LNS) [Shaw, 1998]

### No improvement

- We move towards a solution in the neighborhood **without improving** the objective  
emph  $\Rightarrow$  Don't be a goldfish

## Tabu Search [Glover, 1986]

### Principle

- We start from a solution  $s$ .
- We move towards **the best** solution in the neighbourhood which is not **forbidden**
- Add  $s$  to the forbidden solutions for the next  $m$  iterations

### Memory

- Prohibiting solutions can be memory-intensive
- Instead we forbid movements

### Aspiration criterion

A tabu movement can be accepted if it leads to a **better** solution than the best solution known so far

## Size of tabu list

- If  $m$  too small, **intensification** too strong  $\Rightarrow$  blocking search around a local optimum
- If  $m$  too large, **diversification** too strong  $\Rightarrow$  risk of missing solutions

### Optimal list length varies

- from one problem to another
- from one instance to another of the same problem
- during the resolution of the same instance

[Battiti, Protasi 2001]: adapt this length dynamically

- Need for diversification  $\Rightarrow$  increase  $m$
- Need for intensification  $\Rightarrow$  decrease  $m$

## Local search

### Principle

- We start with an initial solution
- At each step, we modify the solution
  - trying to improve the value of the objective function
  - in the hope of obtaining the global optimum
- Local approach
  - depending on the problem no guarantee of optimality (heuristic)
  - low-cost

### Note

- This assumes the existence of an objective function
- What if there isn't one?

## Constraint Based Local Search

### Principle

- Given a problem of the form
  - $\mathcal{V} = \{v_1, \dots, v_n\}$ : variables
  - $\mathcal{D} = \{D_1, \dots, D_n\}$ : domains
  - $\mathcal{C} = \{C_1, \dots, C_p\}$ : constraints
- Objective function to minimize: number of unsatisfied constraints

### Intuition

- Search guided by problem structure
  - constraints give structure to the problem and variables link them together
- any type of constraint can be used

## Constraint Based Local Search

### N-queens

- on a  $n \times n$  chessboard
- Place  $n$  queens so that no queen can capture another one

### Formulation

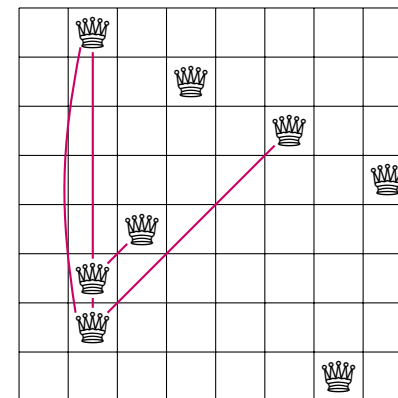
- $l_j$ : queen's column on line  $i$
- $l_i \neq l_j$
- $l_i + i \neq l_j + j$  (upward diagonal)
- $l_i - i \neq l_j - j$  (downward diagonal)

### Objective function

- Number of unsatisfied constraints

## Constraint Based Local Search

### Example

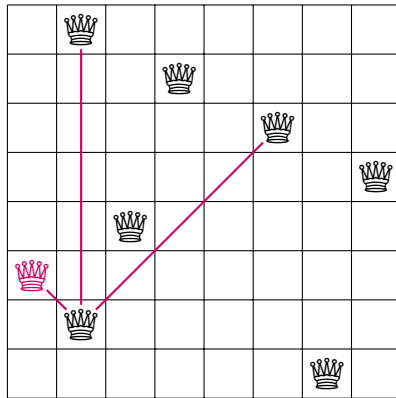


Objective function: 5



## Constraint Based Local Search

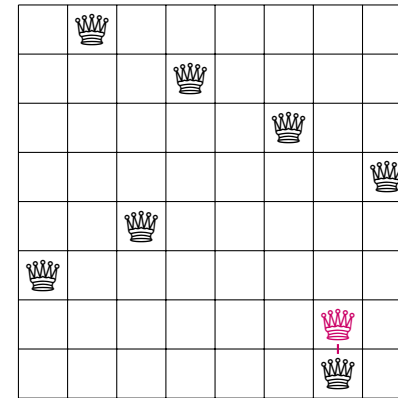
Example



Objective function: 3

## Constraint Based Local Search

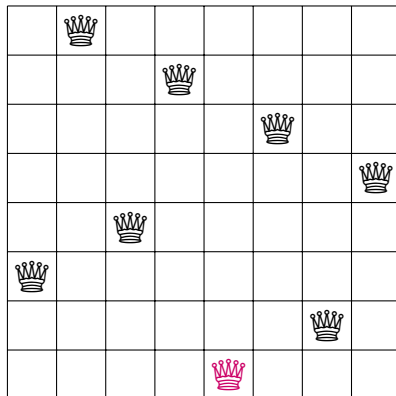
Example



Objective function: 1

## Constraint Based Local Search

Example



Objective function: 0